

On successively orthogonal systems of operations

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Systems of k -ary operations generalizing orthogonal sets are considered.

These systems have the following property: every k successive k -ary operations, $k \geq 2$, of the system are orthogonal.

We call these systems successively orthogonal, establish some properties, give examples and methods of construction of these systems.

A k -ary operation A (briefly, a k -operation) on a set Q is a mapping $A : Q^k \rightarrow Q$ defined by $A(x_1^k) \rightarrow x_{k+1}$, and in this case write $A(x_1^k) = x_{k+1}$.

A k -groupoid (Q, A) is a set Q with one k -ary operation A , defined on Q .

The k -operation $E_i : E_i(x_1^k) = x_i$, $1 \leq i \leq k$, on Q is called *the i -th identity operation (or the i -th selector) of arity k* .

An i -invertible k -operation A , defined on Q , is a k -operation with the following property: the equation $A(a_1^{i-1}, x, a_{i+1}^k) = a_{k+1}$ has a unique solution for each fixed k -tuple $(a_1^{i-1}, a_{i+1}^k, a_{k+1})$ of Q^k .

All 2-invertible binary operations, given on a set Q , form the group $(\Lambda_2; \cdot)$ under the multiplication $(A \cdot B)(x, y) = A(x, B(x, y))$.

A k -ary quasigroup (or simply a k -quasigroup) is a k -groupoid (Q, A) such that the k -operation A is i -invertible for each $i = 1, 2, \dots, k$.

Definition 1 [1]. A k -tuple $\langle A_1, A_2, \dots, A_k \rangle = \langle A_1^k \rangle$ of k -operations, given on a set Q , is called *orthogonal* if the system $\{A_i(x_1^k) = a_i\}_{i=1}^k$ has a unique solution for all $a_i^k \in Q^k$.

A k -tuple $\langle A_1^k \rangle$ of k -operations is orthogonal if and only if the mapping $\theta = (A_1^k) : Q^k \rightarrow Q^k$, $(x_1^k) \rightarrow (A_1(x_1^k), A_2(x_1^k), \dots, A_k(x_1^k)) = (A_1^k)(x_1^k)$ is a permutation on Q^k [1].

Definition 2 [1]. A set $\{A_1, A_2, \dots, A_t\}$, $t \geq k$, of k -operations is called *orthogonal* if every k -tuple of these k -operations is orthogonal.

Definition 3 [1]. A set $\Sigma = \{A_1^t\}$, $t \geq 1$, of k -ary operations, given on a set Q , is called *strongly orthogonal* if the set $\bar{\Sigma} = \{A_1^t, E_1^k\}$ is orthogonal.

Definition 4. A system $\Sigma = \{A_1^t\}$, $t \geq k$, of k -ary operations, given on a set Q , $|Q| \geq 3$, is called *successively orthogonal system (briefly, a SOS)* if any successive k operations are orthogonal.

Every orthogonal set of k -operations is a successively orthogonal system.

Let (Q, A) be a quasigroup, $A^i(x, y) = A(x, A^{i-1}(x, y))$, $i = 2, \dots$.

Theorem 1. If A, A_1, A_2, \dots, A_t are binary quasigroups of the order s_0, s_1, \dots, s_t respectively, in the group $(\Lambda_2; \cdot)$ of all 2-invertible binary operations, given on a set Q , then the sequence

$$F, E, A, A^2, \dots, A^{s_0-1}, F, E, A_1, A_1^2, \dots, A_1^{s_1-1}, \\ F, E, A_2, A_2^2, \dots, A_2^{s_2-1}, \dots, F, E, A_t, A_t^2, \dots, A_t^{s_t-1}$$

is a SOS.

Proposition 1. Let $\Sigma_1 = \{A_1, A_2, \dots, A_{s_1}\}$, $\Sigma_2 = \{B_1, B_2, \dots, B_{s_2}\}$ be strongly orthogonal sets of k -operations. Then the system

$$\Sigma_3 = \{E_1, E_2, \dots, E_k, A_1, A_2, \dots, A_{s_1}, E_1, E_2, \dots, E_k, B_1, B_2, \dots, B_{s_2}\}$$

is a SOS.

Theorem 2. Let A be an 1-invertible k -operation on a set Q , $\theta = (E_2, E_3, \dots, E_k, A) = (E_2^k, A)$, and s_0 be the order of the permutation θ in the group S_{Q^k} , then the system of k -operations

$$\begin{aligned} &E_1, E_2, \dots, E_k, A, A\theta, A\theta^2, \dots, A\theta^{k-1}, A\theta^k, \dots, A\theta^{s_0-k-1}, \\ &E_1, E_2, \dots, E_k, A, A\theta, A\theta^2, \dots, A\theta^{k-1}, A\theta^k, \dots, A\theta^{s_0-k-1}, \dots \end{aligned}$$

is successively orthogonal.

Corollary 1. In the theorem 2 the k -operation $A\theta^{s_0-k-1}$ is k -invertible.

In [2] for a function $f: Q^k \rightarrow Q$ it was defined a complete k -recursive code $K(n/f^{(0)}, f^{(1)}, \dots, f^{(n-k-1)})$ with the check functions: $f^{(0)} = f, f^{(1)}, \dots, f^{(n-k-1)}$. The function $f^{(i)}$ is called the i -th recursive derivative of a function f and is defined recursively as follows:

$$\begin{aligned} f^{(0)}(x_1^k) &= f(x_1^k), \quad f^{(1)}(x_1^k) = f(x_2^k, f^{(0)}(x_1^k)), \dots, \\ f^{(i)}(x_1^k) &= f(x_{i+1}^k, f^{(0)}(x_1^k), \dots, f^{(i-1)}(x_1^k)) \text{ for } i < k, \text{ and} \\ f^{(i)}(x_1^k) &= f(f^{(i-k)}(x_1^k), f^{(i-k+1)}(x_1^k), \dots, f^{(i-1)}(x_1^k)) \text{ for } i \geq k. \end{aligned}$$

V. Izbash and P. Syrbu in [3, Proposition 2] proved that if a k -operation f is a k -quasigroup, then $f^{(i)} = f\theta^i, i = 1, 2, \dots$, where

$$\theta: Q^k \rightarrow Q^k, \theta(x_1^k) = (x_2, x_3, \dots, x_k, f(x_1^k))$$

for all $(x_1^k) \in Q^k$.

A k -quasigroup operation f ($k \geq 2$) is called *recursively r -differentiable* if all its k -recursive derivatives $f^{(0)}, f^{(1)}, \dots, f^{(r)}$ are k -quasigroups [2].

A k -quasigroup we call *strongly recursively r -differentiable* if it is recursively r -differentiable and $r = s_0 - k - 1$, where s_0 is the order of the permutation $\theta = (E_2^k, A)$. In this case $A^{(r+1)} = E_1$. For the binary case this notion was introduced in [4]. From Theorem 2 we obtain the following corollary for any 1-invertible k -function f .

Corollary 2. If f is an 1-invertible k -function, then $f^{(i)} = f\theta^i, i = 1, 2, \dots$, where $\theta = (E_2^k, f)$.

The sequence of the recursive derivatives has the form $E_1, E_2, \dots, E_k, f, f\theta, f\theta^2, \dots, f\theta^{k-1}, f\theta^k, \dots, f\theta^{s_0-k-1}, E_1, E_2, \dots, E_k, f, f\theta, f\theta^2, \dots, f\theta^{k-1}, f\theta^k, \dots, f\theta^{s_0-k-1}, \dots$, where s_0 is the order of the permutation θ .

If f is an r -differentiable k -quasigroup, then $r \leq s_0 - k - 1$.

If a k -quasigroup is strongly recursively r -differentiable, then $r = s_0 - k - 1$.

For an 1-differentiable k -quasigroup $s_0 \geq k + 2$.

Theorem 3. Let a permutation (E_2^k, A) have the order s_0 , then a successively orthogonal system of Theorem 2 contains s_0 different k -operations.

If $s_0 = k + 1$, then the k -operation A is a quasigroup k -operation.

For any 1-invertible k -operation $s_0 \geq k + 1$.

Theorem 4. Let A, A_1, \dots, A_t be 1-invertible k -operations and the permutations $\theta = (E_2^k, A), \theta_1 = (E_2^k, A_1), \dots, \theta_t = (E_2^k, A_t)$ have the order s_0, s_1, \dots, s_t respectively, then the system

$$\begin{aligned} & E_1, E_2, \dots, E_k, A, A\theta, A\theta^2, \dots, A\theta^{k-1}, A\theta^k, \dots, A\theta^{s_0-k-1}, \\ & E_1, E_2, \dots, E_k, A_1, A_1\theta_1, A_1\theta_1^2, \dots, A_1\theta_1^{k-1}, A_1\theta_1^k, \dots, A_1\theta_1^{s_1-k-1}, \dots, \\ & E_1, E_2, \dots, E_k, A_t, A_t\theta_t, A_t\theta_t^2, \dots, A_t\theta_t^{k-1}, A_t\theta_t^k, \dots, A_t\theta_t^{s_t-k-1} \end{aligned}$$

is a SOS.

Proposition 2. Any orthogonal set of k -operations can be continued to a SOS.

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