# On recursively differentiable $k$-quasigroups 

Parascovia Syrbu, Elena Cuzneţov


#### Abstract

Recursive differentiability of linear $k$-quasigroups ( $k \geq 2$ ) is studied in the present work. A $k$-quasigroup is recursively $r$-differentiable ( r is a natural number) if its recursive derivatives of order up to $r$ are quasigroup operations. We give necessary and sufficient conditions of recursive 1-differentiability (respectively, $r$-differentiability) of the $k$-group $(Q, B)$, where $B\left(x_{1}, \ldots, x_{k}\right)=x_{1} \cdot x_{2} \cdot \ldots$. $x_{k}, \forall x_{1}, x_{2}, \ldots, x_{k} \in Q$, and ( $Q, \cdot \cdot$ ) is a finite binary group (respectively, a finite abelian binary group). The second result is a generalization of a known criterion of recursive $r$-differentiability of finite binary abelian groups [4]. Also we consider a method of construction of recursively $r$-differentiable finite binary quasigroups of high order $r$. The maximum known values of the parameter $r$ for binary quasigroups of order up to 200 are presented.


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The notions "recursive derivative" and "recursively differentiable quasigroup" were introduced in [1], where the authors considered recursive MDS-codes (Maximum Distance Separable codes). The recursive derivative of order $t \geq 0$ of a $k$-ary groupoid $(Q, A)$ is denoted by $A^{(t)}$ and is defined as follows:

$$
\begin{aligned}
& A^{(0)}=A \\
& A^{(t)}\left(x_{1}^{k}\right)=A\left(x_{t+1}, \ldots, x_{k}, A^{(0)}\left(x_{1}^{k}\right), \ldots, A^{(t-1)}\left(x_{1}^{k}\right)\right) \text { if } 1 \leq t<k \\
& A^{(t)}\left(x_{1}^{k}\right)=A\left(A^{(t-k)}\left(x_{1}^{k}\right), \ldots, A^{(t-1)}\left(x_{1}^{k}\right)\right) \text { if } t \geq k, \forall x_{1}, \ldots, x_{k} \in Q,
\end{aligned}
$$

where we denoted the sequence $x_{1}, x_{2}, \ldots, x_{k}$ by $x_{1}^{k}$. A $k$-ary quasigroup $(Q, A)$ is called recursively $r$-differentiable if the recursive derivatives $A^{(0)}, A^{(1)}, \ldots, A^{(r)}$ are quasigroup operations ( $r \geq 0$ ).

The length $n$ of the codewords in a $k$-recursive code

$$
C(n, A)=\left\{\left(x_{1}, \ldots, x_{k}, A^{(0)}\left(x_{1}^{k}\right), \ldots, A^{(n-k-1)}\left(x_{1}^{k}\right)\right) \mid x_{1}, \ldots, x_{k} \in Q\right\}
$$

given on an alphabet $Q$ of $q$ elements, where $A: Q^{k} \rightarrow Q$ is the defining $k$-ary quasigroup operation, satisfies the condition $n \leq r+k+1$, where $r$ is the maximum order of recursive differentiability of $(Q, A)$. On the other hand, $C(n, A)$ is an MDScode if and only if $d=n-k+1$, where $d$ is the minimum Hamming distance of this code. At present it is an open problem to determine all triplets $(n, d, q)$ of natural numbers such that there exists an MDS-code $C$ of lenght $n$, on an alphabet of $q$ elements, with $|C|=q^{k}$ and with the minimum Hamming distance $d$, for each

[^0]$k \geq 2$. This general problem implies, in particular, the problem of determining the maximum order of recursive differentiability of finite $k$-ary quasigroups $(k \geq 2)$.

Let $(Q, *)$ be a binary quasigroup. Denoting by ${ }_{*}^{*}$ the recursive derivative of order $t$ of the operation $*$, we have:

$$
\begin{gathered}
x * y=x * y \\
x * y=y *(x * y), \\
x * y=\left(x^{t-2} * y\right) *\left(x^{t-1} * y\right), \forall t \geq 2 \text { and } \forall x, y \in Q .
\end{gathered}
$$

It is known that there exist recursively 1-differentiable binary finite quasigroups of any order, except $1,2,6$, and possibly $14,18,26$ and 42 [1]. Some estimations of the maximum (known) order $r$ of recursive differentiability of finite $n$-quasigroups $(n \geq 2)$ are given in [1-4]. General properties of recursively differentiable binary quasigroups are studied in $[4,6,7]$.

The recursive differentiability of $k$-ary quasigroups is closely connected to the orthogonality of the recursive derivatives $[1,4,6]$. It is shown in [1] that a $k$-quasigroup defines an MDS-code of length $n$ if and only if its first $n-k-1$ recursive derivatives are strongly orthogonal. Hence the defining $k$-quasigroup operation of a recursive MDS-code of length $n$ is recursively ( $n-k-1$ )-differentiable. On the other hand, it is known that a system of binary quasigroups is strongly orthogonal if and only if it is (simply) orthogonal [5]. Another "special property" of binary quasigroups is given in [1]: the recursive derivatives of order up to $r$ of a finite binary quasigroup $(Q, *)$ are quasigroup operations if and only if $(Q, *)$ defines a recursive MDS-code of length $r+3$. So, a finite binary quasigroup $(Q, *)$ is recursively $r$-differentiable if and only if its recursive derivatives of order up to $r$ are mutually orthogonal. The last statement implies the fact that there do not exist recursively 1-differentiable quasigroups of orders 2 and 6 and that $r \leq q-2$, where $q=|Q|$ and $r$ is the order of the recursive differentiability of the quasigroup $Q$. Recall that there do not exist orthogonal latin squares of order 2 or 6 , and the number of mutually orthogonal latin squares on a set of $q$ elements does not exceed $q-1$ [5]. The mentioned above results imply the following lemma.

Lemma 1. The maximum order $r$ of recursive differentiability of a finite binary quasigroup of order $q$ satisfies the inequality $r \leq q-2$.

It is shown in [1] that there exist recursively ( $q-2$ )-differentiable finite binary quasigroups of every primary order $q \geq 3$. However, it is an open problem to find the maximum order $r$ of recursive differentiability of finite $k$-ary quasigroups of order $q$, for $k \geq 2$ and non-primary $q$.

Recursive differentiability of linear $n$-ary quasigroups ( $n \geq 2$ ) is studied in the present work. In particular, we give necessary and sufficient conditions of recursive 1-differentiability (respectively, $r$-differentiability) of an $n$-group $(Q, B)$, where $B\left(x_{1}^{n}\right)=x_{1} \cdot x_{2} \cdot \ldots \cdot x_{n}, \forall x_{1}, x_{2}, \ldots, x_{n} \in Q$, and $(Q, \cdot)$ is a finite binary group (respectively, a finite abelian binary group). The second result is a generalization of a known criterion for finite binary abelian groups, given in [4]. Also we consider a
method of construction of recursively differentiable finite binary quasigroups of high (in particular, maximum) order $r$. The maximum known values of the order $r$ of recursive differentiability of finite binary quasigroups of order up to 200, are qiven in Table 1.

Lemma 2. Let $n \geq 2$ and let $\left(Q_{i}, A_{i}\right)$ be a recursively $r_{i}$-differentiable $n$-quasigroup, $i=1, \ldots, m$. Then the direct product $\left(Q_{1} \times \ldots \times Q_{m}, B\right)$,

$$
\begin{equation*}
B\left(\left(x_{11}^{1 m}\right), \ldots,\left(x_{n 1}^{n m}\right)\right)=\left(A_{1}\left(x_{11}^{n 1}\right), \ldots, A_{m}\left(x_{1 m}^{n m}\right)\right), \tag{1}
\end{equation*}
$$

$\forall\left(x_{11}^{1 m}\right), \ldots,\left(x_{n 1}^{n m}\right) \in Q_{1} \times \ldots \times Q_{m}$, is a recursively $r$-differentiable $n$-quasigroup, where $r=\min \left\{r_{1}, \ldots, r_{m}\right\}$.

Proof. Remind that an $n$-ary groupoid $(Q, B)$ is an $n$-ary quasigroup if each element $u_{i}$ in the equality $B\left(u_{1}, \ldots, u_{n}\right)=u_{n+1}$ is uniquely determined by the remaining $n$ elements. Hence, we get from (1) that $\left(Q_{1} \times \ldots \times Q_{m}, B\right)$ is an $n$-quasigroup. To find the recursive derivatives of $B$ we'll consider the following two cases:
(i) $1 \leq t<n$
$B^{(t)}\left(\left(x_{11}^{1 m}\right), \ldots,\left(x_{n 1}^{n m}\right)\right)=$
$=B\left(\left(x_{t+1,1}^{t+1, m}\right), \ldots,\left(x_{n 1}^{n m}\right), B^{(0)}\left(\left(x_{11}^{1 m}\right), \ldots,\left(x_{n 1}^{n m}\right)\right), \ldots, B^{(t-1)}\left(\left(x_{11}^{1 m}\right), \ldots,\left(x_{n 1}^{n m}\right)\right)\right)=$
$B\left(\left(x_{t+1,1}^{t+1, m}\right), \ldots,\left(x_{n 1}^{n m}\right),\left(A_{1}^{(0)}\left(x_{11}^{n 1}\right), \ldots, A_{m}^{(0)}\left(x_{1 m}^{n m}\right)\right), \ldots,\left(A_{1}^{(t-1)}\left(x_{11}^{n 1}\right), \ldots, A_{m}^{(t-1)}\left(x_{1 m}^{n m}\right)\right)\right)=$ $=\left(A_{1}^{(t)}\left(x_{11}^{n 1}\right), \ldots, A_{m}^{(t)}\left(x_{1 m}^{n m}\right)\right) ;$
(ii) $t \geq n$
$B^{(t)}\left(\left(x_{11}^{1 m}\right), \ldots,\left(x_{n 1}^{n m}\right)\right)=B\left(B^{(t-n)}\left(\left(x_{11}^{1 m}\right), \ldots,\left(x_{n 1}^{n m}\right)\right), \ldots, B^{(t-1)}\left(\left(x_{11}^{1 m}\right), \ldots,\left(x_{n 1}^{n m}\right)\right)\right)=$
$=B\left(\left(A_{1}^{(t-n)}\left(x_{11}^{n 1}\right), \ldots, A_{m}^{(t-n)}\left(x_{1 m}^{n m}\right)\right), \ldots,\left(A^{(t-1)}\left(x_{11}^{n 1}\right), \ldots, A_{m}^{(t-1)}\left(x_{1 m}^{n m}\right)\right)\right)=$
$=\left(A_{1}\left(A_{1}^{(t-n)}\left(x_{11}^{n 1}\right), \ldots, A_{1}^{(t-1)}\left(x_{11}^{n 1}\right)\right), \ldots, A_{m}\left(A_{m}^{(t-n)}\left(x_{1 m}^{n m}\right), \ldots, A_{m}^{(t-1)}\left(x_{1 m}^{n m}\right)\right)\right)=$
$=\left(A_{1}^{(t)}\left(x_{11}^{n 1}\right), \ldots, A_{m}^{(t)}\left(x_{1 m}^{n m}\right)\right)$.
Hence, $B^{(t)}\left(\left(x_{11}^{1 m}\right), \ldots,\left(x_{n 1}^{n m}\right)\right)=\left(A_{1}^{(t)}\left(x_{11}^{n 1}\right), \ldots, A_{m}^{(t)}\left(x_{1 m}^{n m}\right)\right)$, for every $t \geq 1$ and every $\left(x_{11}^{1 m}\right), \ldots,\left(x_{n 1}^{n m}\right) \in Q_{1} \times \ldots \times Q_{m}$. As each of the operations $A_{1}, \ldots, A_{m}$ is recursively $r$-differentiable, where $r=\min \left\{r_{1}, \ldots, r_{m}\right\}$, we get that $B$ is recursively $r$-differentiable.

Proposition 1. Let $(Q, \cdot)$ be a finite binary group and $n \geq 2$. The $n$-ary group $(Q, B)$, where $B\left(x_{1}^{n}\right)=x_{1} \cdot x_{2} \cdot \ldots \cdot x_{n}, \forall x_{1}, x_{2}, \ldots, x_{n} \in Q$, is recursively 1-differentiable if and only if the mapping $x \rightarrow x^{2}$ is a bijection in $(Q, \cdot)$.

Proof. The $n$-group $(Q, B)$ is recursively 1-differentiable if and only if the recursive derivative $B^{(1)}$ is a quasigroup operation, i.e. if and only if in the equality

$$
\begin{equation*}
B^{(1)}\left(x_{1}, \ldots, x_{n}\right)=b, \tag{2}
\end{equation*}
$$

every $n$ elements uniquely determine the remaining ( $n+1$ )-th one. Taking $x_{j}=a_{j} \in$ $Q$ in (2), for every $j=2, \ldots, n$, we get the equation $x_{1} \cdot a_{2} \cdot \ldots \cdot a_{n}=b$, which has
a unique solution in $Q$. For $i \in\{2, \ldots, n\}$, taking $x_{j}=a_{j} \in Q, \forall j \neq i, j \in\{1, \ldots, n\}$, we have:

$$
\begin{gathered}
B^{(1)}\left(a_{1}, \ldots, a_{i-1}, x_{i}, a_{i+1}, \ldots, a_{n}\right)=b \Leftrightarrow \\
\Leftrightarrow B\left(a_{2}, \ldots, a_{i-1}, x_{i}, a_{i+1}, \ldots, a_{n}, B\left(a_{1}, \ldots, a_{i-1}, x_{i}, a_{i+1}, \ldots, a_{n}\right)\right)=b \Leftrightarrow \\
\Leftrightarrow a_{2} \cdot \ldots \cdot a_{i-1} \cdot x_{i} \cdot a_{i+1} \cdot \ldots \cdot a_{n} \cdot a_{1} \cdot \ldots \cdot a_{i-1} \cdot x_{i} \cdot a_{i+1} \cdot \ldots \cdot a_{n}=b \Leftrightarrow \\
\Leftrightarrow\left(a_{1} \cdot \ldots \cdot a_{i-1} \cdot x_{i} \cdot a_{i+1} \cdot \ldots \cdot a_{n}\right)^{2}=a_{1} \cdot b
\end{gathered}
$$

Hence, denoting $a_{1} \cdot \ldots \cdot a_{i-1} \cdot x_{i} \cdot a_{i+1} \cdot \ldots \cdot a_{n}$ by $y$, we get that the $n$-group $(Q, B)$ is recursively 1 -differentiable if and only if, for each $b \in Q$, the equation $y^{2}=b$ has a unique solution.

Corollary 1. There exist finite recursively 1-differentiable n-quasigroups of any odd order $q \geq 3$, for every $n \geq 2$.

Proof. This statement follows from the fact that the mapping $x \rightarrow x^{2}$ is a bijection in every finite binary group of odd order $q \geq 3$.

Theorem 1. Let $(Q, \cdot)$ be a finite binary abelian group and let $n \geq 2, r \geq 1$ be two natural numbers. The n-ary group $(Q, B)$, where $B\left(x_{1}^{n}\right)=x_{1} \cdot x_{2} \cdot \ldots \cdot x_{n}$, for every $x_{1}, x_{2}, \ldots, x_{n} \in Q$, is recursively $r$-differentiable if and only if the mappings $x \rightarrow x^{s_{i}^{k}}$ are bijections in the group $(Q, \cdot), \forall i=1, \ldots, n$ and $\forall k=1, \ldots, r$, where the sequences $\left(s_{i}^{k}\right)_{k \geq 0}$ are defined as follows:

1. $k=0$

$$
s_{1}^{0}=\ldots=s_{n}^{0}=1 ;
$$

2. $1 \leq k<n$

$$
\begin{aligned}
& s_{t}^{k}=s_{t}^{0}+\ldots+s_{t}^{k-1}, \quad \forall t=1, \ldots, k \\
& s_{t}^{k}=1+s_{t}^{0}+\ldots+s_{t}^{k+1}, \quad \forall t=k+1, \ldots, n
\end{aligned}
$$

3. $k \geq n$

$$
s_{t}^{k}=s_{t}^{k-n}+\ldots+s_{t}^{k-1}, \forall t=1, \ldots, n .
$$

Proof. As $(Q, \cdot)$ is an abelian group and $B^{(0)}\left(x_{1}^{n}\right)=x_{1} \cdot \ldots \cdot x_{n}$, the recursive derivatives $B^{(1)}$ and $B^{(2)}$ as follows:

$$
\begin{gathered}
B^{(1)}\left(x_{1}^{n}\right)=B\left(x_{2}, \ldots, x_{n}, B^{(0)}\left(x_{1}^{n}\right)\right)=x_{2} \cdot \ldots \cdot x_{n} \cdot x_{1} \cdot \ldots \cdot x_{n}=x_{1} \cdot x_{2}^{2} \cdot \ldots \cdot x_{n}^{2} ; \\
B^{(2)}\left(x_{1}^{n}\right)=B\left(x_{3}, \ldots, x_{n}, B^{(0)}\left(x_{1}^{n}\right), B^{(1)}\left(x_{1}^{n}\right)\right)=x_{3} \cdot \ldots \cdot x_{n} \cdot x_{1} \cdot \ldots \cdot x_{n} \cdot x_{1} \cdot x_{2}^{2} \cdot \ldots \cdot x_{n}^{2}= \\
=x_{1}^{2} \cdot x_{2}^{3} \cdot x_{3}^{4} \cdot \ldots \cdot x_{n}^{4} .
\end{gathered}
$$

Let denote $B^{(k)}\left(x_{1}^{n}\right)=x_{1}^{s_{1}^{k}} \cdot x_{2}^{s_{2}^{k}} \cdot \ldots \cdot x_{n}^{s_{n}^{k}}$, for every $k \geq 0$. To find the sequences $\left(s_{i}^{k}\right)_{k \geq 0}$, where $i=1, \ldots, n$, we will consider the following two cases:

1. $0 \leq k<n$

$$
\begin{aligned}
& B^{(k)}\left(x_{1}^{n}\right)=B\left(x_{k+1}, \ldots, x_{n}, B^{(0)}\left(x_{1}^{n}\right), \ldots, B^{(k-1)}\left(x_{1}^{n}\right)\right)= \\
& =x_{k+1} \cdot \ldots \cdot x_{n} \cdot x_{1}^{s_{1}^{0}} \cdot \ldots \cdot x_{n}^{s_{n}^{0}} \cdot \ldots \cdot x_{1}^{s_{1}^{k-1}} \cdot \ldots \cdot x_{n}^{s_{n}^{k-1}}= \\
& =x_{1}^{s_{1}^{0}+\ldots+s_{1}^{k-1}} \cdot \ldots \cdot x_{k}^{s_{k}^{0}+\ldots+s_{k}^{k-1}} \cdot x_{k+1}^{1+s_{k+1}^{0}+\ldots+s_{k+1}^{k-1}} \cdot \ldots \cdot x_{n}^{1+s_{n}^{0}+\ldots+s_{n}^{k-1}} ;
\end{aligned}
$$

2. $k \geq n$

$$
\begin{aligned}
& B^{(k)}\left(x_{1}^{n}\right)=B\left(B^{(k-n)}\left(x_{1}^{n}\right), \ldots, B^{(k-1)}\left(x_{1}^{n}\right)\right)=B^{(k-n)}\left(x_{1}^{n}\right) \cdot \ldots \cdot B^{(k-1)}\left(x_{1}^{n}\right)= \\
& =x_{1}^{s_{1}^{k-n}+\ldots+s_{1}^{k-1}} \cdot \ldots \cdot x_{n}^{s_{n}^{k-n}+\ldots+s_{n}^{k-1}} .
\end{aligned}
$$

The recursive derivatives $B^{(k)}$, where $k=1,2, \ldots, r$, are quasigroup operations if and only if the mappings $x \rightarrow x^{s_{i}^{k}}$ are bijections in the group $(Q, \cdot), \forall i=1, \ldots, n$ and $\forall k=1, \ldots, r$.

Corollary 2. [4] A finite binary abelian group ( $Q, \cdot$ ) is recursively $r$-differentiable $(r \geq 1)$ if and only if the mappings $x \rightarrow x^{s_{i}^{k}}$ are bijections, $\forall i=1,2$ and $\forall k=1, \ldots, r$, where the sequences $\left(s_{1}^{k}\right)_{k \geq 0}$ and $\left(s_{2}^{k}\right)_{k \geq 0}$ are defined as follows:

$$
s_{1}^{0}=s_{2}^{0}=1 ; \quad s_{1}^{1}=1, s_{2}^{1}=2 ; s_{i}^{k}=s_{i}^{k-2}+s_{i}^{k-1}, \forall k \geq 2, \forall i=1,2 .
$$

Note that $\left(s_{1}^{k}\right)_{k \geq 0}$ and $\left(s_{2}^{k}\right)_{k \geq 0}$ are Fibonacci sequences.
We will give bellow an algorithm of construction of binary linear (over $\mathbb{Z}_{n}$ ) quasigroups, which are recursively differentiable of high order.

Lemma 3. [7] If $(Q, *)$ is a binary quasigroup then, for every $x, y \in Q$ and $\forall s \geq 1$,

$$
\begin{equation*}
x \stackrel{s}{*} y=y^{s-1} *(x * y) . \tag{3}
\end{equation*}
$$

Lemma 4. Let $a \in \mathbb{Z} \backslash\{0\}$ and $x * y=a x+y, \forall x, y \in \mathbb{Z}$. The following statements hold:

1. There exist $u_{s}, v_{s} \in \mathbb{Z}$ such that $x \stackrel{s}{*} y=u_{s} x+v_{s} y, \forall x, y \in \mathbb{Z}, \forall s \geq 1$;
2. If $n \geq 2$ is a natural number, $k \in\{1, \ldots, n-1\}$ and $a=n-k$, then there exists $b_{s+2} \in \mathbb{Z}$ such that $v_{s+2}=n b_{s-2}+\left(-k c_{s}+c_{s+1}\right)$, for $\forall s \geq 1$, where $c_{s}$ and $c_{s+1}$ are the rests from dividing $v_{s}$ and $v_{s+1}$ by $n$, respectively.

Proof. 1. In this case $x \stackrel{1}{*} y=y *(x * y)=a x+(a+1) y, \forall x, y \in \mathbb{Z}$. Denoting $x \stackrel{s-1}{*} y=u_{s-1} x+v_{s-1} y$ and using the mathematical induction and (3), we get $x \stackrel{s}{*} y=u_{s-1} y+v_{s-1}(a x+y)=a v_{s-1} x+\left(u_{s-1}+v_{s-1}\right) y$.
2. As $x_{*}^{s+2} y=(x * y) *\left(x^{s+1} y\right)=\left(a u_{s}+u_{s+1}\right) x+\left(a v_{s}+v_{s+1}\right) y$, the following equalities hold:

$$
\begin{aligned}
v_{s+2}= & a v_{s}+v_{s+1}=(n-k)\left(n b_{s}+c_{s}\right)+\left(n b_{s+1}+c_{s+1}\right)= \\
& =n\left(n b_{s}+c_{s}-k b_{s}+b_{s+1}\right)+\left(-k c_{s}+c_{s+1}\right)
\end{aligned}
$$

where $c_{s}$ and $c_{s+1}$ are the rests from dividing $v_{s}$ and $v_{s+1}$ by $n$, respectively.

Now, let consider the operation $x * y=\bar{a} x+y$ on the ring $\mathbb{Z}_{n}$ of integers modulo $n$, where $(a, n)=1$. Then $\left(\mathbb{Z}_{n}, *\right)$ is a quasigroup and, according to the previous lemma, there exist $\overline{u_{s}}, \overline{v_{s}} \in \mathbb{Z}_{n}$ such that $x \stackrel{s}{*} y=\overline{u_{s}} x+\overline{v_{s}} y, \forall s \geq 0$.

Theorem 2. Let $n \geq 2, a=n-k, k \in\{1, \ldots, n-1\},(a, n)=1$ and $x * y=\bar{a} x+y, \forall x, y \in \mathbb{Z}_{n}$. If, for some $s \geq 1$, the recursive derivatives $\left(\mathbb{Z}_{n}, *\right)$ and $\left(\mathbb{Z}_{n}, \stackrel{s+1}{*}\right)$, where $x \stackrel{i}{*} y=\overline{u_{i}} x+\overline{v_{i}} y, i=s, s+1$, are quasigroups, then $\left(\mathbb{Z}_{n}, \stackrel{s+2}{*}\right)$ is a quasigroup if and only if $\left(-k c_{s}+c_{s+1}, n\right)=1$, where $c_{s}$ and $c_{s+1}$ are the rests from dividing $v_{s}$ and $v_{s+1}$ by $n$, respectively.

Proof. We have: $x{ }_{*}^{s+2} y=\overline{u_{s+2}} x+\overline{v_{s+2}} y=\overline{a v_{s+1}} x+\left(\overline{u_{s+1}}+\overline{v_{s+1}}\right) y$, so $\overline{v_{s+2}}=$ $\overline{u_{s+1}}+\overline{v_{s+1}}=\overline{-k c_{s}+c_{s+1}}$, where $c_{s}$ and $c_{s+1}$ are the rests from dividing $v_{s}$ and $v_{s+1}$ by $n$, respectively. If $\left(\mathbb{Z}_{n}, \stackrel{s}{*}\right)$ and $\left(\mathbb{Z}_{n}, \stackrel{s+1}{*}\right)$ are quasigroups, then $\left(a v_{s+1}, n\right)=1$, hence $\left(\mathbb{Z}_{n}, \stackrel{s+2}{*}\right)$ is a quasigroup if and only if $\left(-k c_{s}+c_{s+1}, n\right)=1$.

Using Theorem 2, we get, for example, that the quasigroups $\left(\mathbb{Z}_{7}, *\right), x * y=4 x+y$, and $\left(\mathbb{Z}_{11}, *\right), x * y=3 x+y$, are recursively 5 - and 9-differentiable, respectively. Recall that the order $r$ of recursive differentiability of a binary quasigroup, defined on a set of $q$ elements, satisfies the inequality $r \leq q-2$. The following corollary gives all values of the element $a$ such that the quasigroup $\left(\mathbb{Z}_{p}, *\right)$, where $x * y=\bar{a} x+y, \forall x, y \in \mathbb{Z}_{p}$, is recursively differentiable of maximum order, for each odd prime $p$, up to 19 .

Corollary 3. Let $\left(\mathbb{Z}_{n}, *\right)$, where $x * y=\bar{a} x+y, \forall x, y \in \mathbb{Z}_{n}$, be a quasigroup. The following statements hold:

1. $\left(\mathbb{Z}_{3}, *\right)$ is recursively 1-differentiable if and only if $a=1$;
2. $\left(\mathbb{Z}_{5}, *\right)$ is recursively 3-differentiable if and only if $a=3$;
3. $\left(\mathbb{Z}_{7}, *\right)$ is recursively 5 -differentiable if and only if $a=1$ or 4 ;
4. $\left(\mathbb{Z}_{11}, *\right)$ is recursively 9-differentiable if and only if $a=3$ or 4;
5. $\left(\mathbb{Z}_{13}, *\right)$ is recursively 11-differentiable if and only if $a=5,8$ or 11 ;
6. $\left(\mathbb{Z}_{17}, *\right)$ is recursively 15 -differentiable if and only if $a=7$ or 10 ;
7. $\left(\mathbb{Z}_{19}, *\right)$ is recursively 17 -differentiable if and only if $a=1,5$ or 7 .

The known estimations $r_{0} \leq r$ of the order $r$ of recursive differentiability of binary finite quasigroups of order $q \leq 200$ are given in the following Table 1. In the cell with coordinates $(m, k)$ we give the known value of the parameter $r$ for quasigroups of order $m+k$. Remark that the cell $(0,0)$ contains the known value of $r$ for the quasigroups of order 200. An analogous table containing the maximum known length of recursive MDS-codes, defined by quasigroups of order up to 100, is given in [2] and we use it in the first ten lines of Table 1.

|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $0(200)$ | $r \geq 2$ | 0 | 0 | 1 | 2 | 3 | 0 | 5 | 6 | 7 |
| 10 | 1 | 9 | 1 | 11 | $?$ | 1 | 14 | 15 | $?$ | 17 |
| 20 | 2 | 2 | 1 | 21 | 2 | 23 | $?$ | 25 | 2 | 27 |
| 30 | 1 | 29 | 30 | 1 | 1 | 3 | 1 | 35 | 1 | 2 |
| 40 | 1 | 39 | $?$ | 41 | 2 | 1 | 1 | 45 | 1 | 47 |
| 50 | 4 | 1 | 2 | 51 | 3 | 3 | 5 | 4 | 4 | 57 |
| 60 | 2 | 59 | 3 | 5 | 62 | 4 | 3 | 65 | 3 | 3 |
| 70 | 4 | 69 | 6 | 71 | 3 | 3 | 3 | 5 | 4 | 77 |
| 80 | 5 | 79 | 3 | 81 | 4 | 4 | 4 | 3 | 6 | 87 |
| 90 | 3 | 5 | 4 | 3 | 4 | 4 | 4 | 95 | 4 | 7 |
| 100 | 2 | 99 | 1 | 101 | 6 | 1 | 1 | 105 | 1 | 107 |
| 110 | 1 | 1 | 5 | 111 | 1 | 3 | 2 | 1 | 1 | 5 |
| 120 | 1 | 119 | 1 | 1 | 2 | 123 | 1 | 125 | 126 | 1 |
| 130 | 1 | 129 | 1 | 5 | 1 | 1 | 6 | 135 | 1 | 137 |
| 140 | 2 | 1 | 1 | 9 | 1 | 3 | 1 | 1 | 2 | 147 |
| 150 | 1 | 149 | 6 | 1 | 1 | 3 | 1 | 155 | 1 | 1 |
| 160 | 3 | 5 | 1 | 161 | 2 | 1 | 1 | 165 | 1 | 167 |
| 170 | 1 | 1 | 2 | 171 | 1 | 3 | 5 | 1 | 1 | 177 |
| 180 | 2 | 179 | 1 | 1 | 6 | 3 | 1 | 9 | 2 | 1 |
| 190 | 1 | 189 | 1 | 191 | 1 | 1 | 2 | 195 | 1 | 197 |

Table 1. Estimations of the parameter $r$ (order of recursive differentiability) in the case of binary quasigroups

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Parascovia Syrbu
Received July 21, 2022
Moldova State University,
Department of Mathematics
E-mail: parascovia.syrbu@gmail.com
Elena Cuzneţov
Moldova State University,
Department of Mathematics
E-mail: lenkacuznetova95@gmail.com


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