## Estimates for the number of vertices with an interval spectrum in proper edge colorings of some graphs

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**Abstract.** For an undirected, simple, finite, connected graph G, we denote by V(G) and E(G) the sets of its vertices and edges, respectively. A function  $\varphi : E(G) \rightarrow \{1, 2, \ldots, t\}$  is called a proper edge *t*-coloring of a graph G if all colors are used and no two adjacent edges receive the same color. An arbitrary nonempty subset of consecutive integers is called an interval. The set of all proper edge *t*-colorings of G is denoted by  $\alpha(G, t)$ . The minimum value of t for which there exists a proper edge *t*-coloring of a graph G is denoted by  $\chi'(G)$ . Let

$$\alpha(G) \equiv \bigcup_{t=\chi'(G)}^{|E(G)|} \alpha(G,t).$$

If G is a graph,  $\varphi \in \alpha(G)$ ,  $x \in V(G)$ , then the set of colors of edges of G incident with x is called a spectrum of the vertex x in the coloring  $\varphi$  of the graph G and is denoted by  $S_G(x,\varphi)$ . If  $\varphi \in \alpha(G)$  and  $x \in V(G)$ , then we say that  $\varphi$  is interval (persistent-interval) for x if  $S_G(x,\varphi)$  is an interval (an interval with 1 as its minimum element). For an arbitrary graph G and any  $\varphi \in \alpha(G)$ , we denote by  $f_{G,i}(\varphi)(f_{G,pi}(\varphi))$ the number of vertices of the graph G for which  $\varphi$  is interval (persistent-interval). For any graph G, let us set

$$\eta_i(G) \equiv \max_{\varphi \in \alpha(G)} f_{G,i}(\varphi), \quad \eta_{pi}(G) \equiv \max_{\varphi \in \alpha(G)} f_{G,pi}(\varphi).$$

For graphs G from some classes of graphs, we obtain lower bounds for the parameters  $\eta_i(G)$  and  $\eta_{pi}(G)$ .

Mathematics subject classification: 05C15.

Keywords and phrases: Proper edge coloring, interval spectrum.

## 1 Introduction

We consider undirected, simple, finite, connected graphs. For a graph G, we denote by V(G) and E(G) the sets of its vertices and edges, respectively. For any  $x \in V(G)$ ,  $d_G(x)$  denotes the degree of the vertex x in G. For a graph G, we denote by  $\Delta(G)$  the maximum degree of a vertex of G. A function  $\varphi : E(G) \to \{1, 2, \ldots, t\}$  is called a proper edge t-coloring of a graph G if all colors are used and no two adjacent edges receive the same color. The set of all proper edge t-colorings of G is denoted by  $\alpha(G, t)$ . The minimum value of t for which there exists a proper edge t-coloring of a graph G is called a chromatic index [22] of G and is denoted by  $\chi'(G)$ .

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Let us also define the set  $\alpha(G)$  of all proper edge colorings of the graph G

$$\alpha(G) \equiv \bigcup_{t=\chi'(G)}^{|E(G)|} \alpha(G,t).$$

If G is a graph,  $\varphi \in \alpha(G)$ ,  $x \in V(G)$ , then the set of colors of edges of G incident with x is called a spectrum of the vertex x in the coloring  $\varphi$  of the graph G and is denoted by  $S_G(x, \varphi)$ .

An arbitrary nonempty subset of consecutive integers is called an interval. An interval with the minimum element p and the maximum element q is denoted by [p, q]. An interval D is called an h-interval if |D| = h.

For any real number  $\xi$ , we denote by  $\lfloor \xi \rfloor$  ( $\lceil \xi \rceil$ ) the maximum (minimum) integer which is less (greater) than or equal to  $\xi$ .

If G is a graph,  $\varphi \in \alpha(G)$ , and  $x \in V(G)$ , then we say that  $\varphi$  is interval (persistent-interval) for x if  $S_G(x,\varphi)$  is a  $d_G(x)$ -interval (a  $d_G(x)$ -interval with 1 as its minimum element). For an arbitrary graph G and any  $\varphi \in \alpha(G)$ , we denote by  $f_{G,i}(\varphi)(f_{G,pi}(\varphi))$  the number of vertices of the graph G for which  $\varphi$  is interval (persistent-interval). For any graph G, let us [17] set

$$\eta_i(G) \equiv \max_{\varphi \in \alpha(G)} f_{G,i}(\varphi), \quad \eta_{pi}(G) \equiv \max_{\varphi \in \alpha(G)} f_{G,pi}(\varphi).$$

The terms and concepts that we do not define can be found in [23].

It is clear that if for any graph  $G \eta_{pi}(G) = |V(G)|$ , then  $\chi'(G) = \Delta(G)$ . For a regular graph G, these two conditions are equivalent:  $\eta_{pi}(G) = |V(G)| \Leftrightarrow \chi'(G) = \Delta(G)$ . It is known [15, 19] that for a regular graph G, the problem of deciding whether or not the equation  $\chi'(G) = \Delta(G)$  is true is *NP*-complete. It means that for a regular graph G, the problem of deciding whether or not the equation  $\eta_{pi}(G) = |V(G)|$  is true is also *NP*-complete. For any tree G, some necessary and sufficient condition for fulfilment of the equation  $\eta_{pi}(G) = |V(G)|$  was obtained in [8]. In this paper, for an arbitrary regular graph G, we obtain a lower bound for the parameter  $\eta_{pi}(G)$ .

If G is a graph,  $R_0 \subseteq V(G)$ , and the coloring  $\varphi \in \alpha(G)$  is interval (persistent-interval) for any  $x \in R_0$ , then we say that  $\varphi$  is interval (persistent-interval) on  $R_0$ .

 $\varphi \in \alpha(G)$  is called an interval coloring of a graph G if  $\varphi$  is interval on V(G).

We define the set  $\mathfrak{N}$  as the set of all graphs for which there is an interval coloring. Clearly, for any graph  $G, G \in \mathfrak{N}$  if and only if  $\eta_i(G) = |V(G)|$ .

The notion of an interval coloring was introduced in [6]. In [6,7,16] it is shown that if  $G \in \mathfrak{N}$ , then  $\chi'(G) = \Delta(G)$ . For a regular graph G, these two conditions are equivalent:  $G \in \mathfrak{N} \Leftrightarrow \chi'(G) = \Delta(G)$  [6,7,16]. Consequently, for a regular graph G, four conditions are equivalent:  $G \in \mathfrak{N}, \chi'(G) = \Delta(G), \eta_i(G) = |V(G)|,$  $\eta_{pi}(G) = |V(G)|$ . It means that for any regular graph G,

1) the problem of deciding whether G has or not an interval coloring is NP-complete,

2) the problem of deciding whether the equation  $\eta_i(G) = |V(G)|$  is true or not is NP-complete.

In this paper, for an arbitrary regular graph G, we obtain a lower bound for the parameter  $\eta_i(G)$ .

We also obtain some results for bipartite graphs. The complexity of the problem of existence of an interval coloring for bipartite graphs is investigated in [3, 9, 21]. In [16] it is shown that for a bipartite graph G with bipartition (X, Y) and  $\Delta(G) = 3$ the problem of existence of a proper edge 3-coloring which is persistent-interval on  $X \cup Y$  (or even only on Y [6, 16]) is NP-complete.

Suppose that G is an arbitrary bipartite graph with bipartition (X, Y) [3]. Then  $\eta_i(G) \ge \max\{|X|, |Y|\}.$ 

Suppose that G is a bipartite graph with bipartition (X, Y) for which there exists a coloring  $\varphi \in \alpha(G)$  persistent-interval on Y. Then  $\eta_{pi}(G) \ge 1 + |Y|$ .

Some attention is paid to  $(\alpha, \beta)$ -biregular bipartite graphs [4, 13, 14, 18] in the case when  $|\alpha - \beta| = 1$ .

We show that if G is a (k-1,k)-biregular bipartite graph,  $k \ge 4$ , then

$$\eta_i(G) \ge \frac{k-1}{2k-1} \cdot |V(G)| + \left\lceil \frac{k}{\left\lceil \frac{k}{2} \right\rceil \cdot (2k-1)} \cdot |V(G)| \right\rceil$$

We show that if G is a (k-1,k)-biregular bipartite graph,  $k \ge 3$ , then

$$\eta_{pi}(G) \ge \frac{k}{2k-1} \cdot |V(G)|.$$

## 2 Results

**Theorem 1** (see [17]). If G is a regular graph with  $\chi'(G) = 1 + \Delta(G)$ , then

$$\eta_{pi}(G) \ge \left\lceil \frac{|V(G)|}{1 + \Delta(G)} \right\rceil.$$

*Proof.* Suppose that  $\beta \in \alpha(G, 1 + \Delta(G))$ . For any  $j \in [1, 1 + \Delta(G)]$ , define

$$V_{G,\beta,j} \equiv \{ x \in V(G) / j \notin S_G(x,\beta) \}.$$

For arbitrary integers j', j'', where  $1 \le j' < j'' \le 1 + \Delta(G)$ , we have

$$V_{G,\beta,j'} \cap V_{G,\beta,j''} = \varnothing$$
 and  $\bigcup_{j=1}^{1+\Delta(G)} V_{G,\beta,j} = V(G).$ 

Hence, there exists  $j_0 \in [1, 1 + \Delta(G)]$  for which

$$|V_{G,\beta,j_0}| \ge \left| \frac{|V(G)|}{1 + \Delta(G)} \right|.$$

Set  $R_0 \equiv V_{G,\beta,j_0}$ . *Case 1.*  $j_0 = 1 + \Delta(G)$ . Clearly,  $\beta$  is persistent-interval on  $R_0$ . *Case 2.*  $j_0 \in [1, \Delta(G)]$ . Define a function  $\varphi : E(G) \to [1, 1 + \Delta(G)]$ . For any  $e \in E(G)$ , set:

$$\varphi(e) \equiv \begin{cases} \beta(e) & \text{if } \beta(e) \notin \{j_0, 1 + \Delta(G)\}, \\ j_0 & \text{if } \beta(e) = 1 + \Delta(G), \\ 1 + \Delta(G) & \text{if } \beta(e) = j_0. \end{cases}$$

It is not difficult to see that  $\varphi \in \alpha(G, 1 + \Delta(G))$  and  $\varphi$  is persistent-interval on  $R_0$ .

**Corollary 1** (see [17]). If G is a cubic graph, then there exists a coloring from  $\alpha(G, \chi'(G))$  which is persistent-interval for at least  $\left\lceil \frac{|V(G)|}{4} \right\rceil$  vertices of G.

**Theorem 2** (see [17]). If G is a regular graph with  $\chi'(G) = 1 + \Delta(G)$ , then

$$\eta_i(G) \ge \left\lceil \frac{|V(G)|}{\left\lceil \frac{1+\Delta(G)}{2} \right\rceil} \right\rceil.$$

*Proof.* Suppose that  $\beta \in \alpha(G, 1 + \Delta(G))$ . For any  $j \in [1, 1 + \Delta(G)]$ , define

 $V_{G,\beta,j} \equiv \{ x \in V(G) / j \notin S_G(x,\beta) \}.$ 

For arbitrary integers j', j'', where  $1 \le j' < j'' \le 1 + \Delta(G)$ , we have

$$V_{G,\beta,j'} \cap V_{G,\beta,j''} = \varnothing$$
 and  $\bigcup_{j=1}^{1+\Delta(G)} V_{G,\beta,j} = V(G).$ 

For any  $i \in [1, \lceil \frac{1+\Delta(G)}{2} \rceil]$ , let us define the subset  $V(G, \beta, i)$  of the set V(G) as follows:

$$V(G,\beta,i) \equiv \begin{cases} V_{G,\beta,2i-1} \cup V_{G,\beta,2i} & \text{if } \Delta(G) \text{ is odd and } i \in [1,\frac{1+\Delta(G)}{2}] \\ & \text{or } \Delta(G) \text{ is even and } i \in [1,\frac{\Delta(G)}{2}], \\ V_{G,\beta,1+\Delta(G)} & \text{if } \Delta(G) \text{ is even and } i = 1 + \frac{\Delta(G)}{2}. \end{cases}$$

For arbitrary integers i', i'', where  $1 \le i' < i'' \le \left\lceil \frac{1+\Delta(G)}{2} \right\rceil$ , we have

$$V(G,\beta,i') \cap V(G,\beta,i'') = \varnothing$$
 and  $\bigcup_{i=1}^{\left\lfloor \frac{1+\Delta(G)}{2} \right\rfloor} V(G,\beta,i) = V(G).$ 

Hence, there exists  $i_0 \in \left[1, \left\lceil \frac{1+\Delta(G)}{2} \right\rceil\right]$  for which

$$|V(G,\beta,i_0)| \ge \left\lceil \frac{|V(G)|}{\left\lceil \frac{1+\Delta(G)}{2} \right\rceil} \right\rceil.$$

Set  $R_0 \equiv V(G, \beta, i_0)$ . *Case 1.*  $i_0 = \left\lceil \frac{1+\Delta(G)}{2} \right\rceil$ . *Case 1.a.*  $\Delta(G)$  is even. Clearly,  $\beta$  is interval on  $R_0$ . *Case 1.b.*  $\Delta(G)$  is odd. Define a function  $\varphi : E(G) \rightarrow [1, 1 + \Delta(G)]$ . For any  $e \in E(G)$ , set:

$$\varphi(e) \equiv \begin{cases} (\beta(e) + 1) (\mod(1 + \Delta(G))) & \text{if } \beta(e) \neq \Delta(G), \\ 1 + \Delta(G) & \text{if } \beta(e) = \Delta(G). \end{cases}$$

It is not difficult to see that  $\varphi \in \alpha(G, 1 + \Delta(G))$  and  $\varphi$  is interval on  $R_0$ . *Case 2.*  $1 \leq i_0 \leq \left\lceil \frac{\Delta(G)-1}{2} \right\rceil$ . Define a function  $\varphi : E(G) \to [1, 1 + \Delta(G)]$ . For any  $e \in E(G)$ , set:

$$\varphi(e) \equiv \begin{cases} (\beta(e) + 2 + \Delta(G) - 2i_0) (\operatorname{mod}(1 + \Delta(G))) & \text{if } \beta(e) \neq 2i_0 - 1, \\ 1 + \Delta(G) & \text{if } \beta(e) = 2i_0 - 1. \end{cases}$$

It is not difficult to see that  $\varphi \in \alpha(G, 1 + \Delta(G))$  and  $\varphi$  is interval on  $R_0$ .

**Corollary 2** (see [17]). If G is a cubic graph, then there exists a coloring from  $\alpha(G, \chi'(G))$  which is interval for at least  $\frac{|V(G)|}{2}$  vertices of G.

**Theorem 3** (see [6,7,16]). Let G be a bipartite graph with bipartition (X, Y). Then there exists a coloring  $\varphi \in \alpha(G, |E(G)|)$  which is interval on X.

**Corollary 3.** Let G be a bipartite graph with bipartition (X, Y). Then  $\eta_i(G) \ge \max\{|X|, |Y|\}$ .

**Theorem 4** (see [1,6,7]). Let G be a bipartite graph with bipartition (X, Y) where  $d_G(x) \leq d_G(y)$  for each edge  $(x, y) \in E(G)$  with  $x \in X$  and  $y \in Y$ . Then there exists a coloring  $\varphi_0 \in \alpha(G, \Delta(G))$  which is persistent-interval on Y.

**Theorem 5.** Suppose G is a bipartite graph with bipartition (X, Y), and there exists a coloring  $\varphi_0 \in \alpha(G, \Delta(G))$  which is persistent-interval on Y. Then, for an arbitrary vertex  $x_0 \in X$ , there exists  $\psi \in \alpha(G, \Delta(G))$  which is persistent-interval on  $\{x_0\} \cup Y$ .

Proof. Case 1.  $S_G(x_0, \varphi_0) = [1, d_G(x_0)]$ . In this case  $\psi$  is  $\varphi_0$ . Case 2.  $S_G(x_0, \varphi_0) \neq [1, d_G(x_0)]$ .

Clearly,  $[1, d_G(x_0)] \setminus S_G(x_0, \varphi_0) \neq \emptyset$ ,  $S_G(x_0, \varphi_0) \setminus [1, d_G(x_0)] \neq \emptyset$ . Since  $|S_G(x_0, \varphi_0)| = |[1, d_G(x_0)]| = d_G(x_0)$ , there exists  $\nu_0 \in [1, d_G(x_0)]$  satisfying the condition  $|[1, d_G(x_0)] \setminus S_G(x_0, \varphi_0)| = |S_G(x_0, \varphi_0) \setminus [1, d_G(x_0)]| = \nu_0$ .

Now let us construct the sequence  $\Theta_0, \Theta_1, \ldots, \Theta_{\nu_0}$  of proper edge  $\Delta(G)$ -colorings of the graph G, where for any  $i \in [0, \nu_0]$ ,  $\Theta_i$  is persistent-interval on Y.

Set  $\Theta_0 \equiv \varphi_0$ .

Suppose that for some  $k \in [0, \nu_0 - 1]$ , the subsequence  $\Theta_0, \Theta_1, \ldots, \Theta_k$  is already constructed.

Let

$$t_k \equiv \max(S_G(x_0, \Theta_k) \setminus [1, d_G(x_0)]),$$
  
$$s_k \equiv \min([1, d_G(x_0)] \setminus S_G(x_0, \Theta_k)).$$

Clearly,  $t_k > s_k$ . Consider the path P(k) in the graph G of maximum length with the initial vertex  $x_0$  whose edges are alternatively colored by the colors  $t_k$  and  $s_k$ . Let  $\Theta_{k+1}$  be obtained from  $\Theta_k$  by interchanging the two colors  $t_k$  and  $s_k$  along P(k).

It is not difficult to see that  $\Theta_{\nu_0}$  is persistent-interval on  $\{x_0\} \cup Y$ . Set  $\psi \equiv \Theta_{\nu_0}$ .

**Corollary 4.** Let G be a bipartite graph with bipartition (X, Y) where  $d_G(x) \leq d_G(y)$ for each edge  $(x, y) \in E(G)$  with  $x \in X$  and  $y \in Y$ . Let  $x_0$  be an arbitrary vertex of X. Then there exists a coloring  $\varphi_0 \in \alpha(G, \Delta(G))$  which is persistent-interval on  $\{x_0\} \cup Y$ .

**Corollary 5** (see [17]). Let G be a bipartite graph with bipartition (X, Y) where  $d_G(x) \leq d_G(y)$  for each edge  $(x, y) \in E(G)$  with  $x \in X$  and  $y \in Y$ . Then  $\eta_{pi}(G) \geq 1 + |Y|$ .

Remark 1. Notice that the complete bipartite graph  $K_{n+1,n}$  for an arbitrary positive integer n satisfies the conditions of Corollary 5. Is is not difficult to see that  $\eta_{pi}(K_{n+1,n}) = 1 + n$ . It means that the bound obtained in Corollary 5 is sharp since in this case |Y| = n.

Remark 2. Let G be a bipartite (k - 1, k)-biregular graph with bipartition (X, Y), where  $k \ge 3$ . Assume that all vertices in X have the degree k - 1 and all vertices in Y have the degree k. Then the numbers  $\frac{|X|}{k}$ ,  $\frac{|Y|}{k-1}$ , and  $\frac{|V(G)|}{2k-1}$  are integer. It follows from the equalities gcd(k - 1, k) = 1 and  $|E(G)| = |X| \cdot (k - 1) = |Y| \cdot k$ .

**Theorem 6** (see [17]). Let G be a bipartite (k - 1, k)-biregular graph, where  $k \ge 4$ . Then

$$\eta_i(G) \ge \frac{k-1}{2k-1} \cdot |V(G)| + \left| \frac{k}{\left\lceil \frac{k}{2} \right\rceil \cdot (2k-1)} \cdot |V(G)| \right|.$$

*Proof.* Suppose that (X, Y) is a bipartition of G. Without loss of generality we assume that all vertices in X have the degree k - 1 and all vertices in Y have the degree k. Clearly,  $\chi'(G) = \Delta(G) = k$ . Suppose that  $\beta \in \alpha(G, k)$ . For any  $j \in [1, k]$ , define:

$$V_{G,\beta,j} \equiv \{ x \in X/j \notin S_G(x,\beta) \}.$$

For arbitrary integers j', j'', where  $1 \le j' < j'' \le k$ , we have

$$V_{G,\beta,j'} \cap V_{G,\beta,j''} = \emptyset$$
 and  $\bigcup_{j=1}^k V_{G,\beta,j} = X.$ 

For any  $i \in [1, \lceil \frac{k}{2} \rceil]$ , let us define the subset  $V(G, \beta, i)$  of the set X as follows:

$$V(G,\beta,i) \equiv \begin{cases} V_{G,\beta,2i-1} \cup V_{G,\beta,2i} & \text{if } k \text{ is odd and } i \in [1,\frac{k-1}{2}] \\ & \text{or } k \text{ is even and } i \in [1,\frac{k}{2}], \\ V_{G,\beta,k} & \text{if } k \text{ is odd and } i = \frac{1+k}{2}. \end{cases}$$

For arbitrary integers i', i'', where  $1 \le i' < i'' \le \left\lceil \frac{k}{2} \right\rceil$ , we have

$$V(G, \beta, i') \cap V(G, \beta, i'') = \emptyset$$
 and  $\bigcup_{i=1}^{\left\lceil \frac{k}{2} \right\rceil} V(G, \beta, i) = X.$ 

Hence, there exists  $i_0 \in \left[1, \left\lceil \frac{k}{2} \right\rceil\right]$  for which

$$|V(G,\beta,i_0)| \ge \left\lceil \frac{|X|}{\left\lceil \frac{k}{2} \right\rceil} \right\rceil.$$

Set  $R_0 \equiv Y \cup V(G, \beta, i_0)$ . It is not difficult to verify that

$$|R_0| \ge \frac{k-1}{2k-1} \cdot |V(G)| + \left\lceil \frac{k}{\left\lceil \frac{k}{2} \right\rceil \cdot (2k-1)} \cdot |V(G)| \right\rceil.$$

Case 1.  $i_0 = \lceil \frac{k}{2} \rceil$ . Case 1.a. k is odd. Clearly,  $\beta$  is interval on  $R_0$ . Case 1.b. k is even. Define a function  $\varphi : E(G) \to [1, k]$ . For any  $e \in E(G)$ , set:

$$\varphi(e) \equiv \begin{cases} (\beta(e) + 1) \pmod{k} & \text{if } \beta(e) \neq k - 1, \\ k & \text{if } \beta(e) = k - 1. \end{cases}$$

It is not difficult to see that  $\varphi \in \alpha(G, k)$  and  $\varphi$  is interval on  $R_0$ . *Case 2.*  $i_0 \in [1, \lceil \frac{k}{2} \rceil - 1]$ . Define a function  $\varphi : E(G) \to [1, k]$ . For any  $e \in E(G)$ , set:

$$\varphi(e) \equiv \begin{cases} (\beta(e) + 1 + k - 2i_0) (\operatorname{mod} k) & \text{if } \beta(e) \neq 2i_0 - 1, \\ k & \text{if } \beta(e) = 2i_0 - 1. \end{cases}$$

It is not difficult to see that  $\varphi \in \alpha(G, k)$  and  $\varphi$  is interval on  $R_0$ .

**Corollary 6** (see [17]). Let G be a bipartite (k - 1, k)-biregular graph, where k is even and  $k \ge 4$ . Then

$$\eta_i(G) \ge \frac{k+1}{2k-1} \cdot |V(G)|.$$

 $\Box$ 

**Corollary 7** (see [17]). Let G be a bipartite (3, 4)-biregular graph. Then there exists a coloring from  $\alpha(G, 4)$  which is interval for at least  $\frac{5}{7}|V(G)|$  vertices of G.

Remark 3. For an arbitrary bipartite graph G with  $\Delta(G) \leq 3$ , there exists an interval coloring of G [10–12]. Consequently, if G is a bipartite (2, 3)-biregular graph, then  $\eta_i(G) = |V(G)|$ .

*Remark* 4. Some sufficient conditions for existence of an interval coloring of a (3, 4)-biregular bipartite graph were obtained in [2, 5, 20].

**Theorem 7** (see [17]). Let G be a bipartite (k-1,k)-biregular graph, where  $k \ge 3$ . Then

$$\eta_{pi}(G) \ge \frac{k}{2k-1} \cdot |V(G)|.$$

*Proof.* Suppose that (X, Y) is a bipartition of G. Without loss of generality we assume that all vertices in X have the degree k - 1 and all vertices in Y have the degree k. Clearly,  $\chi'(G) = \Delta(G) = k$ . Suppose that  $\beta \in \alpha(G, k)$ .

For any  $j \in [1, k]$ , define:

$$V_{G,\beta,j} \equiv \{ x \in X/j \notin S_G(x,\beta) \}.$$

For arbitrary integers j', j'', where  $1 \le j' < j'' \le k$ , we have

$$V_{G,\beta,j'} \cap V_{G,\beta,j''} = \varnothing$$
 and  $\bigcup_{j=1}^k V_{G,\beta,j} = X.$ 

Hence, there exists  $j_0 \in [1, k]$  for which

$$|V_{G,\beta,j_0}| \ge \frac{|X|}{k}.$$

Set  $R_0 \equiv Y \cup V_{G,\beta,j_0}$ . It is not difficult to verify that

$$|R_0| \ge \frac{k}{2k-1} \cdot |V(G)|.$$

Case 1.  $j_0 = k$ . Clearly,  $\beta$  is persistent-interval on  $R_0$ . Case 2.  $j_0 \in [1, k - 1]$ . Define a function  $\varphi : E(G) \to [1, k]$ . For any  $e \in E(G)$ , set:

$$\varphi(e) \equiv \begin{cases} \beta(e) & \text{if } \beta(e) \notin \{j_0, k\}, \\ j_0 & \text{if } \beta(e) = k, \\ k & \text{if } \beta(e) = j_0. \end{cases}$$

It is not difficult to see that  $\varphi \in \alpha(G, k)$  and  $\varphi$  is persistent-interval on  $R_0$ .

**Corollary 8** (see [17]). Let G be a bipartite (3, 4)-biregular graph. Then there exists a coloring from  $\alpha(G, 4)$  which is persistent-interval for at least  $\frac{4}{7}|V(G)|$  vertices of G.

Acknowledgment. The author thanks professors A. S. Asratian and P. A. Petrosyan for their attention to this work.

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Received April 30, 2013